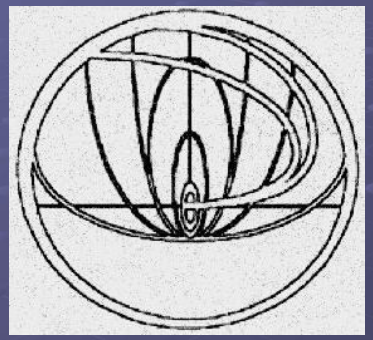
Systems and Devices 2 (Network)
Lec 3a: Transport Layer

Before we get started ...

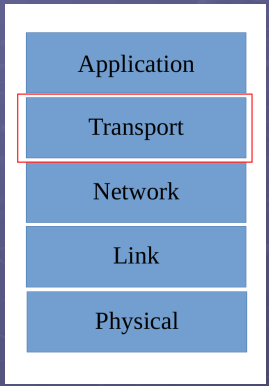
- How did your research into NTP go?
  - RFC 958 (5905): https://tools.ietf.org/html/rfc958
  - Fixed point representation: 32, 64 or 128 bit formats
  - Base date: 0 h 1 January 1900 UTC
    - Coordinated Universal Time (UTC)
  - Dooms day : 8 Feb 2036 (wrap around, reset to 0)
- From the top level view point we have all the protocols we need to get our network started, BUT, what is happening under the bonnet?
  - What are the UDP and TCP protocols we have identified in the transport layer?
  - How do we actually pass data between machines?
  - Why do we need two types of protocols?

Before we get started ...



- John Titor : in 2000 posts appeared on discussion boards claiming to be an American military time traveller from 2036, looking for a IBM 5100 computer, so that he could fix various legacy machines in the future, as it was not susceptible to this problem.

Internet protocol stack

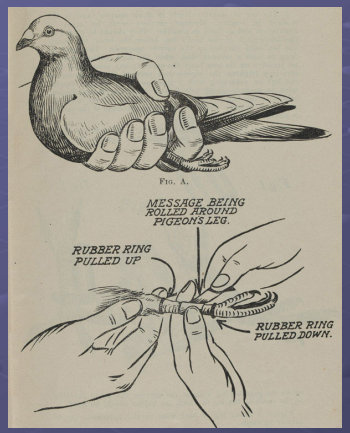


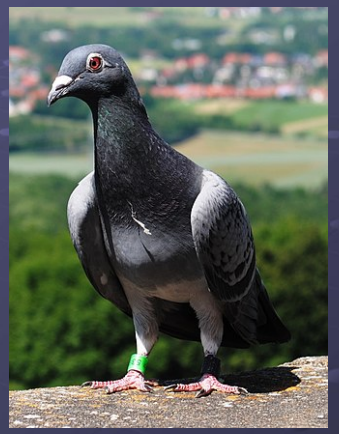
- Application layer
- Transport
  - Breaks messages down into **segments** that will be transferred. Also deals with error detection, congestion control and retransmission in the event of lost packets.
    - Transmission Control Protocol (TCP)
    - User Datagram Protocol (UDP)
- Network
- Link
- Physical

IPoAC

- To help explain these concepts we will consider one of the most important internet protocols ever developed : IPoAC
  - RFC 1149: https://tools.ietf.org/html/rfc1149
  - RFC 2549: https://tools.ietf.org/html/rfc2549
  - RFC 6214: https://tools.ietf.org/html/rfc6214
  - Note, typically updated in April. Especially high data rates in London, located around Trafalgar Square.

IPoAC





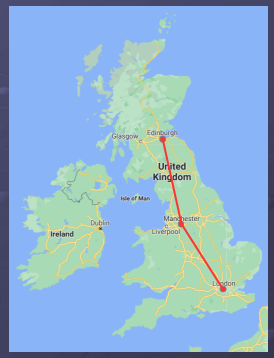
- IP over Avian Carriers
  - Message encapsulated in a metal canister attached to leg

IPoAC

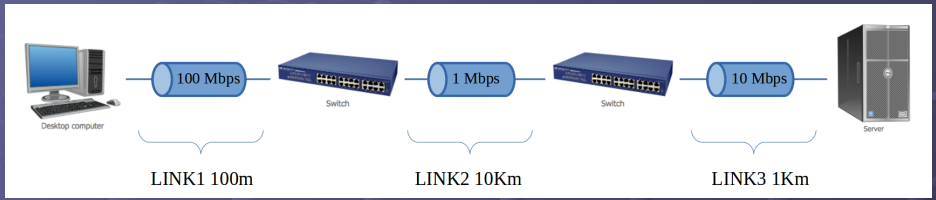
- Multiple packets in flight at the same time, these may be received out of order. Unfortunately packet lose may also occur :(

IPoAC

- Delays associated with “transmitting” packets from London to Edinburgh
  - Processing delay : time taken to load message into canister.
    - Also may need to encrypt message
  - Queuing delay : pigeons that have already arrived from different messages i.e. already in the hutch and waiting to be released.
  - Transmission delay : time taken to move pigeons out of the hutch and “throw” into the air, avoiding mid-air collisions. D = Len / Rate
  - Propagation delay : Time Of Flight (ToF)
- $D_{End\text{-}End} = N_{stages}(D_{proc} + D_{Que} + D_{Tran} + D_{Prop})$



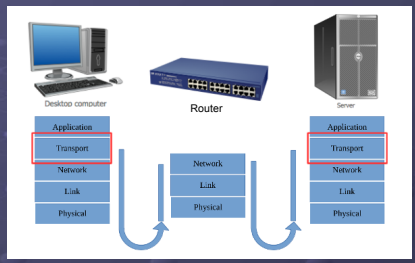
Quick Quizzz



- If we assume that the propagation delay on each link is 2x10^8 m/sec i.e. max is speed of light, but slower in different materials, and that packet length is 1500 bits.
  - What is the transmission and propagation delays of each link
  - What is the end-to-end delay to transfer each packet?
    - You will need to make some assumptions / estimations.

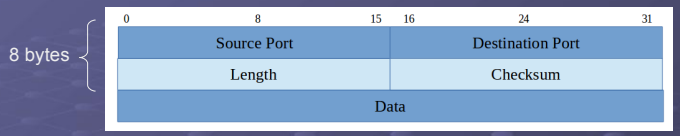
Pause to consider ...

- The transport layer provides a logical communications path between client and server.
  - Not concerned about the network structure, or how data is transmitted e.g. pigeons :)



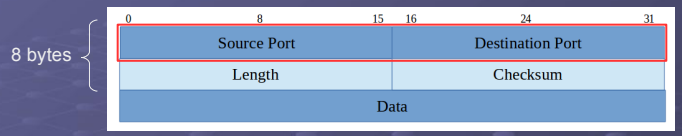
  - Application layer processes can use this end-to-end communications link to pass data (segments) i.e. a logical link between transport layers on “different” machines.
    - Routers / switches do not need to process the transport layer to route packets across the network i.e. a logical communication link between hosts is provided by the network layer.

UDP



- User Datagram Protocol
  - RFC 768 : https://tools.ietf.org/html/rfc768
    - In RFC UDP packets referred to as a datagrams, but this name also used to describe network layer, so to avoid confusion will refer to as segments :)
  - Created in 1983 as a simple, low overhead (small header) communications protocol.
    - Very simple Ports, Length and Checksum.
  - Defined as an unreliable, connectionless service
    - Best-effort delivery service, no handshaking / ACK, as soon as data is ready TX i.e. fire and forget.

UDP



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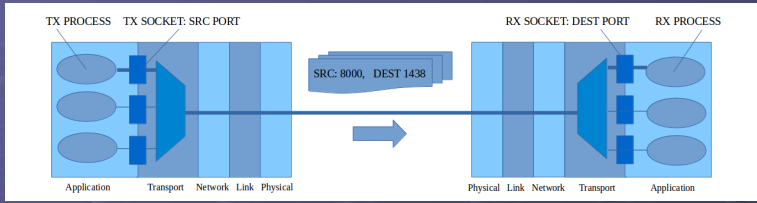
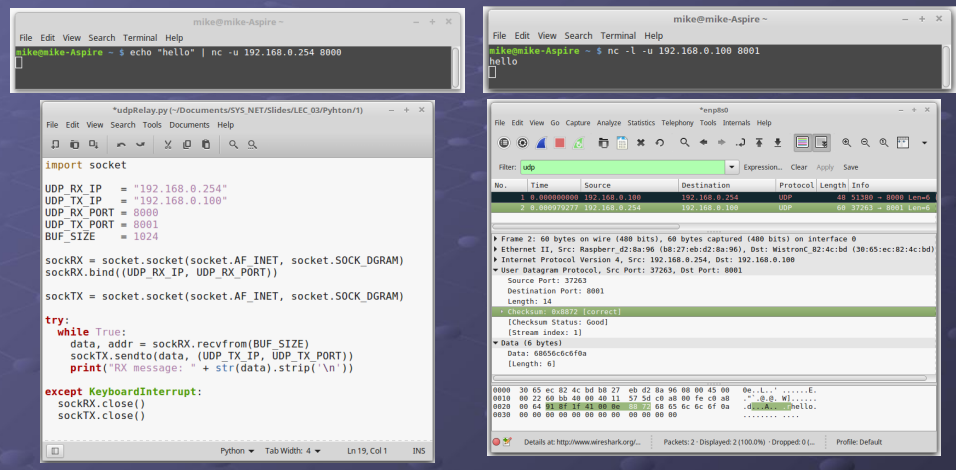


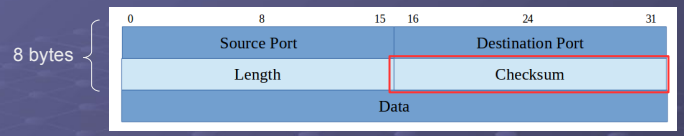
- Q : hosts can have multiple network connections open at a time. How does it decide which process gets received data and how does it identify the process on a remote host to send data to?
  - Port numbers : 16bit port numbers 0 - 65335
    - IANA: https://bit.ly/35mKNR9
    - Categorised as system, registered and ephemeral ports
  - Network layer performs host-to-host data delivery. Process-to-process data delivery is transport-layer **multiplexing** and **demultiplexing** via socket port numbers.

Demo



- Simple python based UDP relay server

UDP



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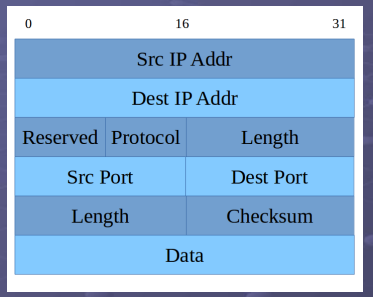
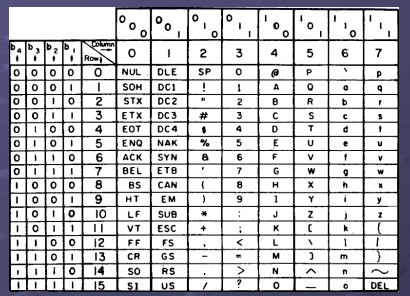
UDP checksum





- To calculate the UDP checksum we use a pseudo header
  - IP addresses included to ensure that the segment received (packet) is for this host and not a corrupted segment.
  - Checksum calculation must be efficient to avoid increasing processing delays. Note, checksum offloading in HW.
    - https://tools.ietf.org/html/rfc1071

UDP checksum

Accumulation	Bit inversion	Check
1111101011110101	00000010 = 0x02	0x0294
+0000011110011110	10010100 = 0x94	+0xFD6B
10000001010010011	= 0x0294	0xFFFF
0000001010010011	11111101 = 0xFD	
+0000000000000001	01101011 = 0x6B	
0000001010010100	= 0xFD6B	

- TX: UDP checksum generated through repeated addition
  - 32bit blocks broken down into 16bit values and accumulated
  - Any overflows wrapped round (bit 16) and added back into the 16bit accumulation, repeated if second overflow
  - Final result inverted and stored in UDP header.

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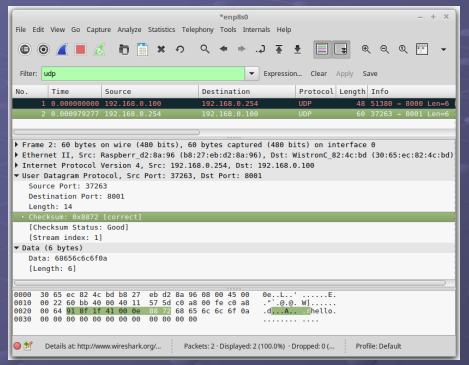
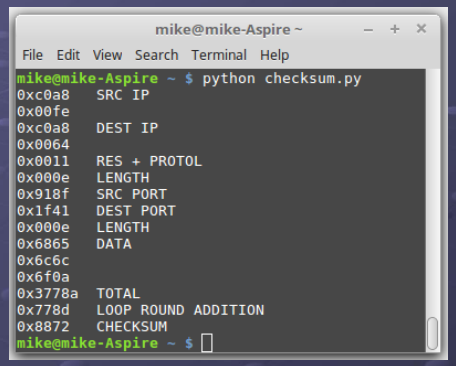
0000001010010100	= 0xFD6B	

- RX: UDP checksum regenerated (removing checksum)
  - Receiving host regenerates checksum and adds received checksum to it. Should generate the result 0xFFFF. If not an error has been detected.

Demo





- UDP checksum generation
  - To simplify the process, performed using a simple python script.

Quick Quizzz

01101101
00101011
+11001100

- Generate the 8bit checksum for the above data.
- If the received checksum is 0x65 has an error occurred?
- Is UDP a connection or connectionless based protocol?
- NTP uses UDP port 123, what happens if a data segment is dropped in the network i.e. what happens if no time data is returned to the client? Does this matter?

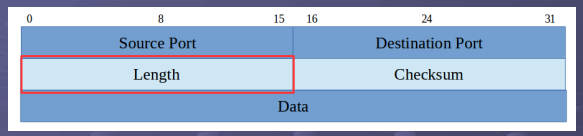
UDP

- Advantages of UDP:
  - Low overheads: small header (8 bytes), no connection establishment delays. Application code can transmit segment as soon as it is ready.
  - Speed: low overheads and data transmitted at full data rate i.e. no congestion control, does not have to wait for ACK.
  - Minimise server loads: do not need to record connection state, send buffers etc. Therefore, server can support more clients.
  - UDP is ideal for real-time applications (less delay) that can tolerant some packet loss.

UDP

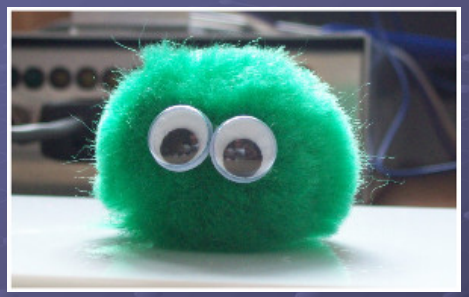
- Disadvantages of UDP
  - No error control, corrupt segments dropped i.e. fail silent. Applications layer must be able to tolerant these scenarios.
    - Can push re-transmit functionality up into the application layer.
  - No congestion control i.e. hosts transmitting large amounts of data can swamp a network i.e. block connections from other machines :(.
  - No confirmation / acknowledgement that the receiving host has received the transmitted data i.e. unreliable, segments could get lost.

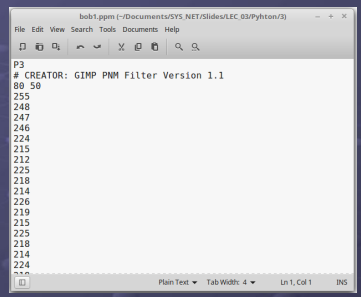
UDP



- Q : what is the biggest possible UDP segment size?
  - RFC 791 states that the max “safe” packet size is 576B
    - “a reasonable sized data block”
    - UDP segment <= 576B : all host must be accept packets of this size i.e. guaranteed to be deliverable, but not guaranteed to be delivered.
    - UDP segment > 576B : are allowed to be dropped by a router / hosts. Limitations in the router / switch hardware
  - Max possible UDP payload dependent on lower layers, but for now we will say approximately 64K - (IP + UDP) headers.
    - UDP packet may be fragmented by lower layers i.e. split up into smaller chunks If any fragment lost the entire packet is dropped.
    - MTU: Maximum Transmission Unit, for IP frames this is ~1500 bytes.

Programming Task





- Write two python programs ie. TX and RX, to transfer a ppm image file of Bob (text file) between PC and Pi.
  - Hint: the file may be larger than the “max” UDP size, break the file down into chunks using the file.read(BUF) method.
- Q : what port should you use? Could you use port 123?

Hints

```
sockTX.sendto(data, (UDP_IP, UDP_PORT))
data = f.read(BUF_SIZE)
```

```
sockRX.settimeout(2)
data,addr = sockRX.recvfrom(BUF_SIZE)
f.write(data)
```

```
sockTX = socket.socket(socket.AF_INET, socket.SOCK_DGRAM)
```

```
UDP_IP = "192.168.101.1"
UDP_PORT = 8000
BUF_SIZE = 1024
```

```
sockRX = socket.socket(socket.AF_INET, socket.SOCK_DGRAM)
sockRX.bind((UDP_IP, UDP_PORT))
```

```
data,addr = sockRX.recvfrom(BUF_SIZE)
number_of_segments = int(data)
```

```
file_size = os.path.getsize(read_file_name)
number_of_segments = (-(-file_size//BUF_SIZE))
```

```
sockTX.sendto(write_file_name.encode("ascii"), (UDP_IP, UDP_PORT))
sockTX.sendto(str(number_of_segments).encode("ascii"), (UDP_IP, UDP_PORT))
```

- What your software needs to do:
  - Send the image file name and the number of segments that will be transmitted i.e. MTU = 1024B.
  - Send the image data as multiple MTU sized segments.
  - Received segments are stored to a file using TX file name.

Summary

- UDP
  - Fast : low overhead (8bytes) communications protocol
    - Note, remember that the UDP header and data will be passed to lower layers in the stack, each adding additional headers e.g. IP header can vary from 20 to 60 bytes.
  - Connectionless : no time wasted setting up a connection between the two machines i.e. as soon as data is ready transmit onto the network and hope the RX host is ready :)
  - Efficient : do not need to reserve memory in server to store connection status or send data buffers i.e. reduces server memory loading.
  - Unreliable : a bit misleading, best effort delivery service.
    - Note, don't confuse this term with faulty, rather that segment delivery is not guaranteed.

Summary

- UDP is used to implement:
  - DNS, NTP, DHCP, BOOTP, TFTP, SNMP, RTP, RTSP, …
  - Steam (game client) ports 27000–27030-ish.
- To overcome the limitations of UDP we need a different protocol:
  - TCP : Transmission Control Protocol
  - For applications were we need to guarantee that data is delivered e.g. email.
  - For next week have a look into this protocol, its a LOT more complex, and do finish off the UDP file transfer coding quizzz.