Systems and Devices 2 (Network) Lec 3b: Transport Layer

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Before we get started ...

- Did you manage to teleport Bob (file transfer)?
 - Could you get my most excellent code working :)
- How did your research into TCP go?
 - Implementing a reliable communications protocol is tricky, need to consider all possible scenarios: lost / corrupt packets, packets received out of order, re-transmission of packets, fair usage of network bandwidth i.e. congestion control ...

• Questions to consider :

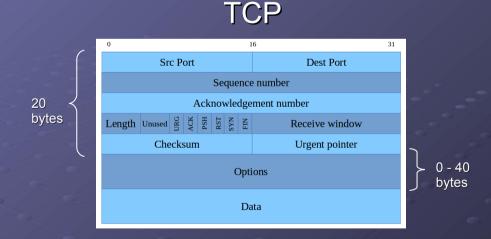
- How does TCP implement a reliable connection?
- What are the differences between UDP and TCP?

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TCP

Transmission Control Protocol

- RFC 793 : https://tools.ietf.org/html/rfc793
- Created in 1981 one of the core internet protocols.
- TCP is a connection orientated protocol i.e. before an application process can send data it must first perform a handshake to ensure that a connection is possible
 - TCP is <u>not</u> an end-to-end protocol like circuit switching, it only existing in the transport layer of the communicating hosts, not in lower layers i.e. no reserved connections, "packets" can take different paths.
- ► TCP uses full duplex connections. If there is a connection between process A and B, then we can send data A→B and B→A at the "same" time.
- TCP implements a reliable, ordered, error-checked data stream between two hosts i.e. retransmission. University of York : M Freeman 2024



• TCP header, depending on option fields used can vary from 20-60B (UDP uses 8B).

Demo

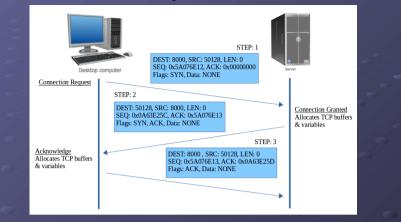
tcpTX.pv (~/Documents/SYS NET/Slides/LEC 03/Pyhton/4) - + ×	tcp.hello.world.pcapng - + ×
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File Edit View Search Tools Documents Help	
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	Piter top
	No. Time Source Destination Protocol Length Info
1 descent conduct	1 0.00000000 037-00 097-00 097-00 097-00 097-00 097-00 097-00 000 000 000 000 000 000 000 000 000
1 import socket	2 0.00059225 192.106.0.254 192.108.0.109 TCP 74 8009 - 50125 193, ACK Sen-0 ACK 1917-55126 197
2 import time	3 0.000735294 192,160,0.100 192,160,0.254 TCP 66 50128 - 0000 [ACK] Sep-1 Ack-1 Min-64256 Len-0 TSys
	4 0.000772926 192.168.0.100 192.168.0.254 TCP 77 50128 - 8090 [P34, ACK] 540+1 ALM-94236 Later
3	5 0.001343116 192.168.0.254 192.168.0.100 TCP 66 8000 - 50128 [ACK] Seg=1 Ack=12 Min=65152 Len=0 TSv
4 TCP IP = "192.168.0.254"	6 0.001905689 192.168.0.254 192.168.0.100 TCP 77 8000 - 50128 [PSH, ACK] Seq=1 Ack=12 Win+65152 Len=
	7 0.001922159 192.168.0.100 192.168.0.254 TCP 66 50128 - 6000 [ACK] Seg=12 Ack=12 Min=64256 Len=0 TS
5 TCP PORT = 8000	8 0.002026001 102.160.0.100 102.168.0.254 TCP 56 50128 - 8000 [FIN, ACK] Seq=12 Ack=12 Win=64256 Len
6 #TCP PORT = 80	9 0.052931764 192.168.0.254 192.108.0.100 TCP 66 3609 - 30128 [FIN, ACK] Seg-12 Ack+12 Min-65132 Len 10.0.652041145 192.168.0.169. 192.168.0.254 TCP 65 36124.0.2018 [Seg-13 Ack-13 Min-65352 Len 10.0.1014 [Seg-13 Ack-13 Min-14535] [Seg-13 Ack-13 Min-14555] [Seg-13 Ack-13 Min-145555] [Seg-13 Ack-13 Min-1455555] [Seg-13 Ack-13 Min-1455555] [Seg-13 Ack-13 Min-14555555555] [Seg
	10 0.00240145 132.160.0.100 132.163.0.254 TCP 66 50126 - 8090 [ACK] Securit Ack-13 Min-64256 Leng-15 11 0.00240145 132.160.0.254 192.168.0.100 TCP 66 8000 - 50126 [ACK] Securit Ack-13 Min-65151 Leng-15
7 BUF SIZE = 1024	TI 0.4053/0/12 125:100.0.524 125:100.0.100 ICA 00.0000 - 20150 DVMI 200412 MINH2025 C6040 I2
	Frame 6: 77 bytes on wire (616 bits), 77 bytes captured (616 bits) on interface 0
0	Frame 6: 77 System on wire (616 0115), 77 System captured (616 0115) on interface 0 Ethernet II, Src: Ramberr distants (616/271ebid2/asi6), Dat: Wisternof 82-66:16(38:65:ec:82:46:16)
9 sockTX = socket.socket(socket.AF INET, socket.SOCK STREAM)	Internet Protoci Version 4. Src: 192.160.0.254 [bt: 192.160.0.100
	* Transmission Control Protocol, Src Port: 0000, Dst Port: 50128, Sep: 1, Ack: 12, Len: 11
10 sockTX.connect((TCP_IP, TCP_PORT))	Source Port: 8000
11	Destination Port: 50128
	[Stream index: 0]
12 try:	[TCP Segment Lon: 11]
13 while True:	Sequence number: 1 (relative sequence number)
	[Next sequence number: 12 (relative sequence number)] Acknowledgment number: 12 (relative ack number)
<pre>14 lowerCase = "hello world"</pre>	1040 Header Length 32 bytes (0)
15 sockTX.send(lowerCase)	Flags: 0x010 (PSH, ACK)
	Nindow size value: 502
<pre>16 upperCase = sockTX.recv(BUF SIZE)</pre>	[Calculated window size: 65152]
	[Window size scaling factor: 128]
	Checksum: Rob354 [unverified]
18 time.sleep(2)	[Checksum Status: Unverified]
	Urgent pointer: 0 > Outions: (12 bytes), No-Operation (NDP), No-Operation (NDP), Timestames
19 except KeyboardInterrupt:	 Options: 112 Options, Ne-Operation (NVP), Ne-Operation (NVP), (Invitanes) SEGUARK analysis]
20 sockTX.close()	Timestanos
	TCP psyload (11 bytes)
21	Data (11 bytes)
22	
	10010 10 65 eC 82 4C bd b8 27 eb 42 8a 96 68 49 45 00 be
23	0020 00 64 1f 40 c3 d0 0a 63 e2 5d 5a 07 6e 1e 00 10d.0c]Z.n
	0010 01 fd b3 54 00 00 01 01 00 0a mc m0 61 cc d1 72
Python Tab Width: 4 Ln 19, Col 1 INS	S M Ready to load or capture Packets: 11 - Displayed: 11 (190.0%) - Dropped: 010.0%) Profile: Default
	Andre A

• TCP server: convert lower case string to upper case.

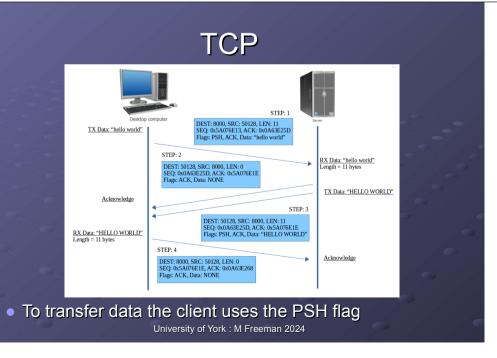
An echo function : **<u>single</u>**, multiple and continuous.

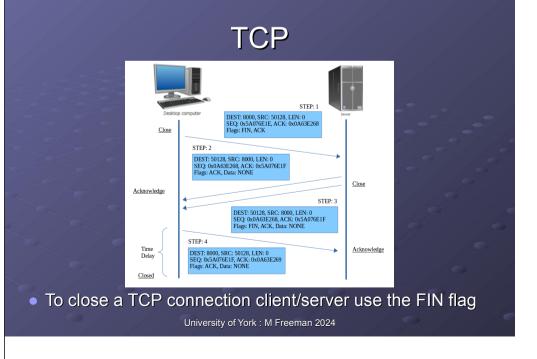
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Three-way handshake



 To initialise a TCP connection the two communicating processes perform a three-way handshake.
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Pause to consider ...

- To ensure reliable connections TCP adds additional initialisation and acknowledgement segments. However, these incur significant overheads i.e. time.
 - To transfer the 10 characters UDP needs 1 segment, TCP needs 9 segments. Therefore, when using TCP we normally keep the link open and transfer multiple values across it.
- This highlights one of the main differences between TCP (stream) vs UDP (message) from the programming point of view i.e. how to separate out different values.
 - UDP: send individual messages i.e. paired send / receive
 - TCP: programmers responsibility to structure the data i.e. either by using a fixed length packets or delimiters such as newline characters.

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Points to note ...

- Reliable data transfer is implemented using positive acknowledgements and timers (discussed next).
- Establishing a TCP link takes time i.e. 3-way handshake. If the client / servers are situated on different continents this will result in significant delays.
- Therefore, web browsers will tend to keep open TCP connection to avoid having to re-establish a link e.g. HTTP communicates across TCP links, using:
 - ▶ Non-Persistent : closed after each object is transferred.
 - Persistent : multiple objects sent over the same link.
 - HTTP1.1 : all connections assumed persistent, however, a web server's default time-out can be 5 – 15 sec.

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Demo tcpTX.py (~/Docum D 🛛 🖌 🖬 🗋 🗙 ဂ 🤇 + + 🤉 ∓ 🛓 🥅 🖬 ୧୯୯୮ Y D B l import socke import time 4 TCP TP = "192 168 0 254" TCP PORT = 8000 6 #TCP POPT = 807 BUF SIZE = 10249 sockTX = socket.socket(socket.AF INET. socket.SOCK STREAM) 10 sockTX.connect((TCP IP, TCP PORT)) 12 try while True: lowerCase = "hello world" sockTX.send(lowerCase) upperCase = sockTX.recv(BUF_SIZE) print("TX:" + upperCase) 18 time.sleep(2) 19 except KeyboardInterrupt: sockTX.close() hon 👻 Tab Width: 4 👻 🛛 In 19 Col 1

TCP server: convert lower case string to upper case.
 An echo function : single, <u>multiple</u> and <u>continuous</u>.

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Programming Task



- Q: reimplement the ppm image file (text file) transfer problem (code on VLE) using TCP, that should be simple :).
 - You will find this tricky. You need to think in terms of a stream of data, rather than packets of data.
 - Consider how the name, number of segments and image data will be sent across the TCP connection. Do you need to send the number of segments?

Quick Quizzz



- The most important parts of the TCP header are the SEQ and ACK numbers i.e. to ensure reliability.
 - SEQ number is the byte stream number of the first byte in the segment.
 - ACK number is the sequence number of the next byte the RX host is expecting from the TX host.
- Q : what are the missing the SEQ and ACK numbers if:
 - RX in order segment 0,1,2 etc, or RX out of order 0,2,1?

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Pause to consider ...



That was the easy bit :), now need to consider the what ifs :(

 IpoAC: multiple packets in flight at the same time, these may be received out of order, packet lose may also occur ... University of York : M Freeman 2024

Time-outs

```
\alpha = 0.125
estRTT = (1 - \alpha) \times estRTT + \alpha \times curRTT
```

 $\beta = 0.25$ devRTT = $(1 - \beta) \times devRTT + \beta \times | curRTT - estRTT$

Timeout = estRTT + 4 × devRTT

• Q: how do we choose the time-out value

- Too short and we will reject valid segments in transit. Too long and we waste time waiting for segments that will never arrive.
- RFC 6298 : round trip time (RTT) delay. Initial time-out set to 1 sec. If time-out occurs, time-out value doubled, reset on next RX segment, do not use times from re-transmitted segments.

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TCP

NextSeqNum ← RAN()

PrevSeqNum - NextSeqNum segment - []

 To help understand how TCP operates consider the following simplified TX pseudo code.

- Three main events
 - Application data
 - Time-out
 - Acknowledgement

while TRUE do
 case EVENT of
 application_layer_data:
 segment.create(data, NextSeqNum)
 if timer.state = stop then
 timer.state - start
 end if
 networkLayer(segment)
 NextSeqNum + len(data)
 time_out:
 networkLayer(segment.select(notACK, smallestSEQ)
 timer.state - start
 rx_ACK:
 if ACK > PrevSeqNum then
 segment.remove(ACK)

```
segment.remove(ACK)
PrevSeqNum - ACK
if segment.NotACK() then
timer.state - start
end if
end case
end while
```

Demo

			pi@pi-1: ~		~ ^ X	pi@pi-2: ~ ✔	~ X
File	Edit	Tabs	Help			File Edit Tabs Help	
i@pi	1:~ \$	echo	"hello" nc	172.1	5.101.9 8000	pi@pi-2:~ \$ sudo tc qdisc add dev eth0 root netem delay 5000ms pi@pi-2:~ \$ nc -l 172.16.101.9 8000 ☐	Ĵ
	_	_		_		test2 popping	arring
File Ed	lit View G	So Capture	Analyze Statistics Te	elephony T	cols Internals Help	- Please check your p	ower su
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Filter:	tcp			*	Expression Clear Apply	ly Save	
		0 172.16.1		TCP		Seq=0 Win=64240 Len=0 MSS=1400 SACK_PERM=1 TSval=1467919412 TSecr=0 WS=128	
		4 172.16.1 1 172.16.1				n] 38920 … 8000 (SYN) Seq=0 Win=64240 Len=0 MSS=1460 SACK_PERM=1 TSval=1467920437 TSecr=0 MS=128 n] 38920 … 8000 [SYN] Seq=0 Win=64240 Len=0 MSS=1460 SACK_PERM=1 TSval=1467922517 TSecr=0 MS=128	
						ACK Sept0 Acks1 Win=65160 Len=0 MSS=1400 SACK PERM=1 TSval=2590721517 TSecr=1467919412 WS=128	
	5.00050984	5 172.16.1	1.9 172.16.101.5	TCP	74 8000 - 38920 [SYN, A		
4	5.00050247	9 172.16.1	01.5 172.16.101.9	TCP	66 38920 - 8000 [ACK] S	Seq=1 Ack=1 Win=64256 Len=0 TSval=1467924412 TSecr=2590721517	
4 5 6	5.00050247	9 172.16.1 0 172.16.1	01.5 172.16.101.9 01.5 172.16.101.9	TCP TCP	66 38920 - 8000 [ACK] S 72 38920 - 8000 [PSH, A	Seq=1 Ack=1 Win=64256 Len=0 TSval=1467924412 TSecr=2590721517 ACK] Seq=1 Ack=1 Win=64256 Len=6 TSval=1467924412 TSecr=2590721517	
4 5 6 7	5.00050247 5.00077359 6.02535657	9 172.16.1 0 172.16.1 1 172.16.1	01.5 172.16.101.9 01.5 172.16.101.9 01.9 172.16.101.5	TCP TCP TCP	66 38928 - 8008 [ACK] S 72 38928 - 8008 [PSH, A 74 [TCP Out-Of-Order] 8	Seq=1 Ack=1 Win=64256 Len=0 TSval=1407924412 TSecr=2500721517 Ack] Seq=1 Ack=4 Win=64256 Len=6 TSval=1467924412 TSecr=2500721517 000 _ 31920 (STW, Ack) Seq=0 Ack=1 Win=6510 Lene0 MSS:400 SACK @EBM=3 TSval=2500722542 TSecr=1467019412 MS=120	
4 5 6 7 8	5.00050247 5.00077359 6.02535657 6.02539410	9 172.16.1 0 172.16.1 1 172.16.1 8 172.16.1	91.5 172.16.101.9 91.5 172.16.101.9 91.9 172.16.101.5 91.5 172.16.101.5 91.5 172.16.101.9	TCP TCP TCP TCP	66 38920 - 8000 [ACK] S 72 38920 - 8000 [PSH, A 74 [TCP Out-Of-Order] 8 66 [TCP Dup ACK 5#1] 38	Sequi Acki Min-64206 Lenn TSW11467924412 TSerr2500712137 ACX Sequi Acki Min-64206 Lenn TSW11467924412 TSerr2500723157 8000 _ 18302 [STN, AcX] Sequi Acki Min-64308 Lenn Missian Lenn Missian Security (Comparison Comparison	
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4 5 6 7 8 9	5.00060247 5.00077359 6.02535657 6.02539410 7.09338194 7.09342895	9 172.16.1 0 172.16.1 1 172.16.1 8 172.16.1 0 172.16.1	01.5 172.16.101.9 01.5 172.16.101.9 01.9 172.16.101.9 01.9 172.16.101.9 01.5 172.16.101.9 01.9 172.16.101.9 01.9 172.16.101.9 01.9 172.16.101.9 01.9 172.16.101.9 01.9 172.16.101.9	TCP TCP TCP TCP TCP	66 38920 - 8000 [ACK] S 72 38920 - 8000 [PSH, A 74 [TCP Dut-07-07der] 8 66 [TCP Dup ACK 5#1] 38 74 [TCP Retransmission] 66 [TCP Dup ACK 5#2] 38	Seq3 Acks Mino6220 Lens Tbw1L30724412 TSer-2309721517 Ack Seq4 Acks Mino6250 Lens Tbw1L30724412 TSer-250972517 MMD _ 3097 [SNR_AS] engl Ack American Lens Resident State State State State State State State State State State State State State State State State State St	
4 5 6 7 8 9 10 11 12	5.00060247 5.00077359 6.02535057 6.02539410 7.09338194 7.09342895 8.10537975 8.10542444	9 172.16.1 0 172.16.1 1 172.16.1 8 172.16.1 0 172.16.1 9 172.16.1 9 172.16.1 9 172.16.1 9 172.16.1	01.5 172.16.101.9 01.5 172.16.101.9 01.9 172.16.101.9 01.9 172.16.101.9 01.9 172.16.101.9 01.9 172.16.101.9 01.9 172.16.101.9 01.9 172.16.101.9 01.9 172.16.101.9 01.9 172.16.101.9 01.9 172.16.101.9	TCP TCP TCP TCP TCP TCP TCP TCP	66 38920 - 8000 [ArX] 5 72 38928 - 8000 [PSH, A 74 [TCP Dut-Of-Order] 8 66 [TCP Dutp ArX 5*1] 38 74 [TCP Rotransmission] 66 [TCP Dup ArX 5*2] 38 74 [TCP Retransmission] 66 [TCP Dup ArX 5*3] 38	Sevil Advi Mindel Bend Toballaffeddin II (2007) Sevil Advi Sevil Advisor (2007) Sevil	
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4 5 6 7 8 9 10 11 12 15 16	5.00060247 5.00077359 6.02535657 6.02539410 7.0938194 8.10537975 8.10542444 10.0010462 12.2852284	9 172.16.1 0 172.16.1 1 172.16.1 8 172.16.1 0 172.16.1 9 172.16.1 9 172.16.1 9 172.16.1 9 172.16.1	$\begin{array}{rrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrrr$	TCP TCP TCP TCP TCP TCP TCP TCP	66 38920 - 8000 [ACK] 5 72 38928 - 8000 [PSH, A 74 [TCP Dut-Of-Order] 0 66 [TCP Dup ACK 587] 38 74 [TCP Dep ACK 587] 38 74 [TCP Dep ACK 587] 38 66 [TCP Dup ACK 587] 38 66 8008 - 38322 [ACK] 5 66 3020 - 8000 [FIN, A	Sevil Advi Mindel Bend Toballaffeddin II (2007) Sevil Advi Sevil Advisor (2007) Sevil	

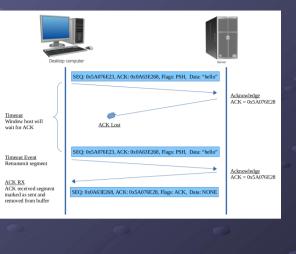
- Command line tool : tc (traffic control)
 - Add 5000ms delay from Pi-1 to Pi-2
 - Work through the time line of errors (black lines), what causes what, tricky :) University of York : M Freeman 2024

Pause to consider ...

• Time-out example 1

 TX host sends a message, it does not receive an ACK, timeout triggered.

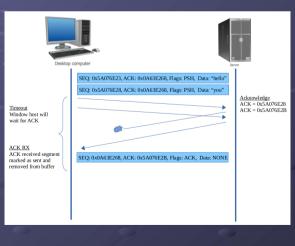
 Q: what happens next i.e. to the second "hello"
 RX by the server?



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Pause to consider

- Time-out example 2
 - TX host sends a message, but only one ACK received.
 - Q: what happens next i.e. does the client need the missing ACK?



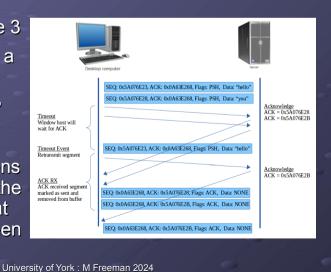
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Pause to consider ...

Time-out example 3

 TX host sends a message, but ACKs delayed, triggering a timeout.

 Q: what happens next i.e. does the client know that all data has been RX?

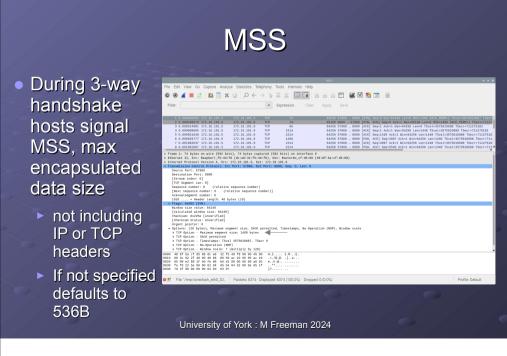


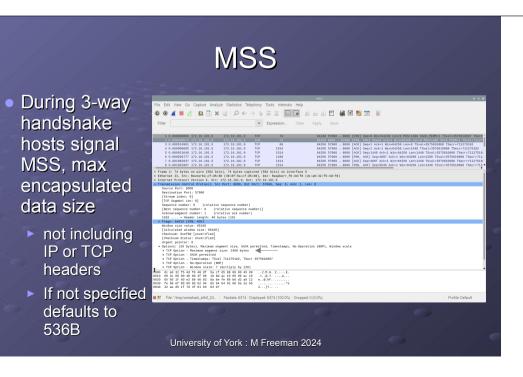
Pause to consider ...

• Time-out example 4 TX host sends a message but one O: 0x5A076E22 ACK: 0x0A62E269 Elage: BSH_Data: of the "words" is SEO: 0x54076E28 ACK: 0x0463E268 Elage: PSH_Data: "an lost SEO: 0x54076E2B_ACK: 0x0A63E268_Elage: PSH_Data: "you ACK TX ACK = 0x5A076F28Q: what happens ACK RX ACK indicates the byte number the host is (Packet Lost) ACK = 0x5A076E28 next? ACK = 0x5A076E28looking for next Timeout Window host will Fast retransmit wait for ACK ×3 ACK RX×3 Same ACK number RX Q: what happens three times triggers fast if only one word is sent? University of York : M Freeman 2024

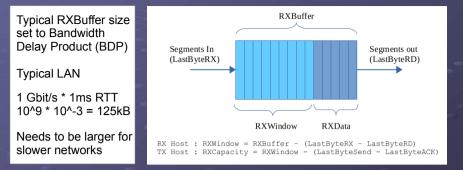
Pause to consider ...

- TX host transmits segments onto transport mechanism. Therefore, there can be multiple segments in flight.
 - $D_{End-End} = N_{stages} (D_{proc} + D_{Que} + D_{Tran} + D_{prop})$
 - Q : if we are transmitting 1000 byte segments to Australia at 1Mbps how many segments can be in flight?
- The RX host receives these segments, but may not process them immediately owing to:
 - Time taken to pass up the layers in the protocol stack
 - The OS is executing other processes.
 - There is a significant processing overhead for the application using these segments.
- Problem : RX host can be swamped with data i.e. data is coming in faster than it can process :(.





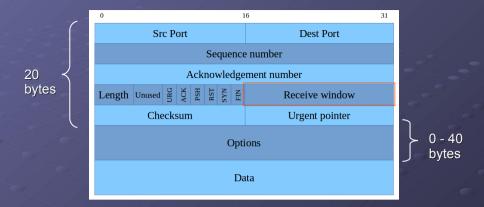
Flow control



 Both TX and RX hosts have buffers to store data when busy. The free space in the RX host's buffer (RXWindow) is transmitted to TX host to ensure buffer overflow does not occur.

- Data is temporarily stored on "wire" i.e. data in flight.
- RX window size transmitted back to host in ACK packets. University of York : M Freeman 2024

Flow control



 Hosts inform each other how much data they can receive using the Receive Window field (16bits)

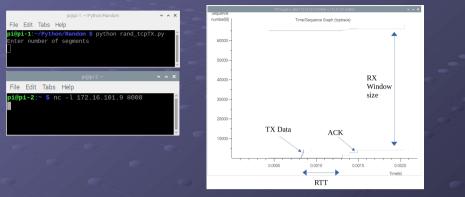
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<text><list-item><list-item>

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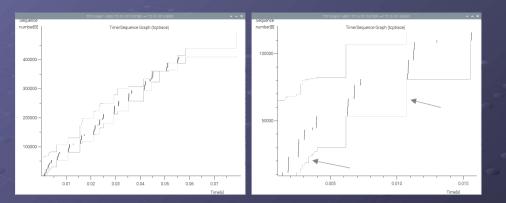
<text><list-item><list-item>

Demo



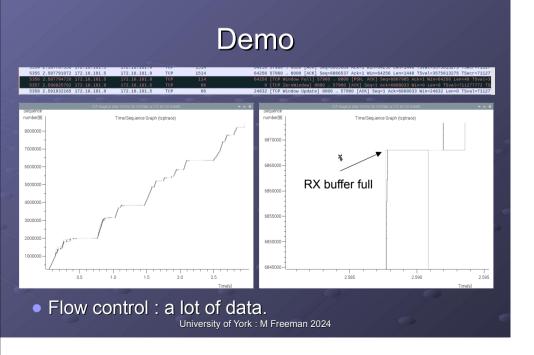
- Flow control : a little data (1).
 - Test code: simple python program transferring one or more TCP packets contain random values. University of York : M Freeman 2024

Demo



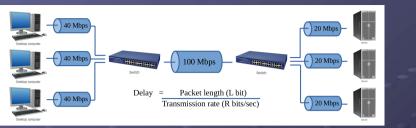
- Flow control : a bit "more" data.
 - More data in flight

Move from individual ACKs to cumulative ACK University of York : M Freeman 2024



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Destination Port: 57960 [Stream index: 0] [TCP Segment Len: 0] Sequence number: 1 (relative [Next sequence number: 1 (rel Acknowledgment number: 6868033 1000 = Header Length: 32 by	<pre>lative sequence number)] (relative ack number)</pre>		1				
Flags: 0x00 (ACK) Windowskie value: 0 [calculated window size: 0] [Window size scaling factor: 122 Checksum 0x040f3 (unverified) [Checksum Status: Unverified] Urgent pointer: 0 0 Options: (12 bytes), No-Operatic	3] on (NOP), No-Operation (NOP), Timestamps						

Pause to consider ...



• We have seen that TCP can adjust TX speed to match the RX speed to avoid buffer overflow.

- However, what should the system do if we just have too much traffic on the network i.e. segments are being lost because of buffer overflows in routers.
 - Remember packet switching uses store and forward. University of York : M Freeman 2024

Congestion control

Congestion Window (CGWindow) = number of unACKed bytes

Number of Bytes in transit $\ \leq$ min(CGWindow, RXWindow) (LastByteSent - LastByteACK) $\ \leq$ min(CGWindow, RXWindow)

Assuming infinite RX buffer, Speed = CGWindow ÷ RTT (bytes per sec)

Initial Speed = MSS ÷ RTT (then increase until missed ACK)

- TCP handles congested networks by using another "window", the Congestion window :). It detects a networks "loading" via the ACK segments
 - ACK received for unACKed segment. All is good increase congestion window size i.e. number of segment in flight.
 - Time-out, or multiple ACKs for same segment RX. Network congested, decrease congestion window size, i.e. packets are being dropped, reduce traffic.

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Congestion control

- TCP congestion-control algorithm
 - RFC 5681: https://tools.ietf.org/html/rfc5681
 - Three core elements: slow start, congestion avoidance and fast recovery (implementations do vary)
- Slow start :
 - Initially CGWnd set to 1 MSS i.e. speed= MSS/RTT.
 - On each RX ACK add 1 MSS to CGWnd i.e. doubled, speed increases exponentially, not too slow :).
 - Continues until ACK time-out i.e. host probes the network to find "max" speed. SSThresold set to CGWnd/2 and CGWnd reset to 1 MSS.
 - Process repeated until SSThresold reached, then switch to congestion avoidance mode.

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Congestion control

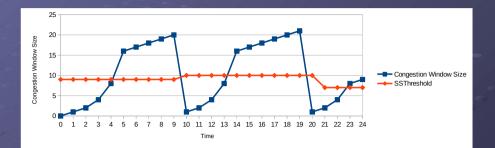
- If three repeated ACKs for the same SEQ number RX, perform fast retransmit and change mode to fast recovery.
- Congestion avoidance:
 - CGWnd is half the value at which congestion was detected. Now rather than doubling CGWnd size it is incremented by 1 MSS when all CGWnd segments have been ACKed i.e. linear increase rather than exponential.
 - If ACK time-out. SSThresold set to CGWnd/2, CGWnd reset to 1 MSS, switch back to soft start mode.
 - If three repeated ACKs for the same SEQ number RX, perform fast retransmit and change mode to fast recovery.

Congestion control

- Fast recovery :
 - CGWnd set to (CGWnd/2)+3MSS. CGWnd increased by MSS for each duplicate ACK RX.
 - If ACK time-out occurs, CGWnd reset to 1 MSS and SSThresold set to CGWnd/2, switch back to soft start mode.
 - If ACK received, CGWnd set to SSThresold, then enter congestion avoidance mode.

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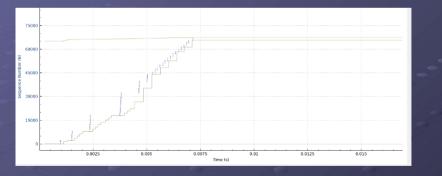
Congestion control



 Note, different strategies taken depending on the version of TCP used. Refer to RFC5681 for additional details e.g. initial CGWnd size, strategies used in each mode to adjust CGWnd size etc.

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Congestion control



 Remember that flow control i.e. RXWindow size, is also active at this time, therefore, speed will also be throttled based on RX host buffer size.

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Summary

- TCP is used to implement:
 - ▶ Telnet, SSH, FTP, HTTP, SMTP, POP ...
 - The internet's workhorse protocol.
- But, to ensure a reliable connection we need additional state information (memory on host), header fields and handshakes in protocol etc, i.e. reliability has a cost.
 - Sequence and acknowledgement numbers to ensure segments are not lost.
 - Time-outs, different types of segments (flags) i.e. SYN, ACK, PSH etc.
 - Initial three way handshake to ensure server is ready to communicate and agree upon parameters used in this transfer e.g. RXWindow sizes.

Summary

- When you look into all the if-buts-and-maybes of how TCP works i.e. all the possible scenarios of time-outs, ACKs etc, its VERY complex.
- We have also not looked at the different option field in the header, or optimisations etc. If you have time you may want to have a read around the area.
- However, we still have some unanswered questions:
 - How does a host know how to route a TCP packet across the Internet i.e. we are not using circuit switching, so how do we know were to send our packets, how do we get from host A to host B?